



# KANGAROOTWELVE

## draft-viguiet-kangarootwelve-01

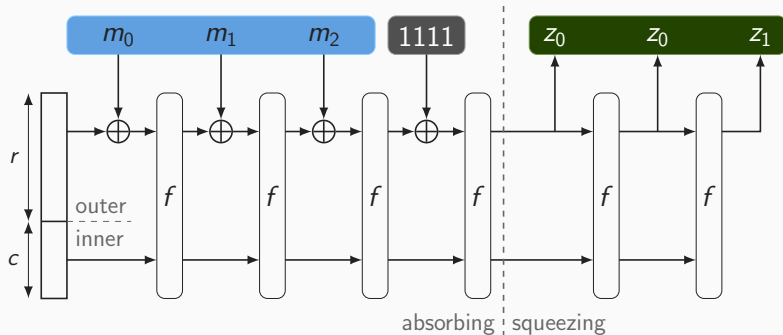
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CFRG Meeting, March 19, 2018

<sup>1</sup>Radboud University, Nijmegen, The Netherlands

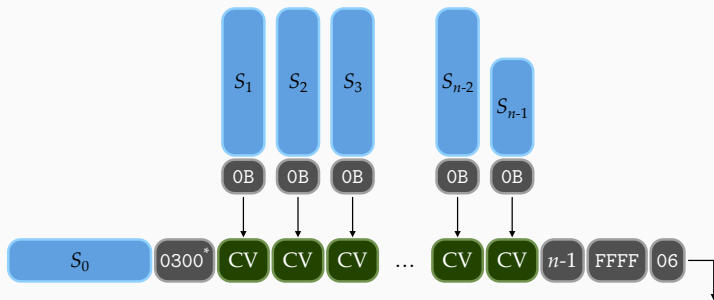
# What is KANGAROOTWELVE?



## ► SHAKE128

- eXtendable Output Function
- Sponge construction
- Uses KECCAK- $p$ [1600,  $n_r = 24$ ]
- BUT no parallelism

# What is KANGAROOTWELVE?



## ► KangarooTwelve

- eXtendable Output Function
- Tree on top of sponge construction
- KECCAK- $p$  reduced from 24 to 12 rounds
- Parallelism grows automatically with input size
- No penalty for short messages

# How secure is KANGAROOTWELVE?

- ▶ Same security claim as SHAKE128: 128 bits of security

- ▶ Sponge generic security

[EuroCrypt 2008] – On the Indifferentiability of the Sponge Construction

- ▶ Parallel mode with proven generic security

[IJIS 2014] – Sufficient conditions for sound tree and sequential hashing modes

[ACNS 2014] – Sakura: A Flexible Coding for Tree Hashing

- ▶ Sponge function on top of KECCAK- $p$ [1600,  $n_T = 12$ ]

- Round function unchanged

⇒ cryptanalysis since 2008 still valid

- Safety margin: from *rock-solid* to *comfortable*

# Status of KECCAK cryptanalysis

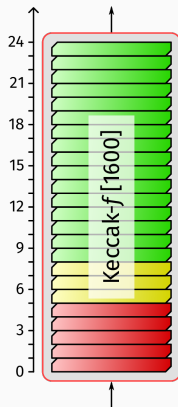
- ▶ Collision attacks up to 5 rounds
  - Also up to 6 rounds, but for non-standard parameters ( $c = 160$ )

[Song, Liao, Guo, CRYPTO 2017]

- ▶ Stream prediction
  - in 8 rounds ( $2^{128}$  time, prob. 1)
  - in 9 rounds ( $2^{256}$  time, prob. 1)

[Dinur, Morawiecki, Pieprzyk, Srebrny, Straus, EUROCRYPT 2015]

- ▶ Lots of third party cryptanalysis available at:  
[https://keccak.team/third\\_party.html](https://keccak.team/third_party.html)



# How fast is KANGAROOTWELVE?

- ▶ At least twice as fast as SHAKE128 on short inputs
- ▶ Much faster when parallelism is exploited on long inputs

	Short input	Long input
Intel® Core™ i5-4570 (Haswell)	3.68 c/b	1.44 c/b
Intel® Core™ i5-6500 (Skylake)	2.89 c/b	1.22 c/b
Intel® Core™ i7-7800X (Skylake-X)	2.35 c/b	0.55 c/b

Single core only.



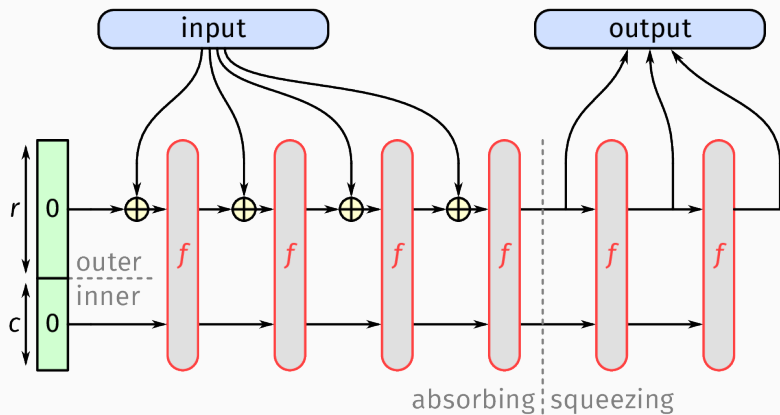
# Why is it interesting for the IETF?

- ▶ KECCAK/KANGAROOTWELVE is an open design
  - Public design rationale
  - Result of an open international competition
  - Long-standing active scrutiny from the crypto community
- ▶ Best security/speed trade-off
  - Speed-up w/o wasting cryptanalysis resources (no tweaks)
  - Proven generic security
- ▶ Scalable parallelism
  - As much parallelism as the implementation can exploit
  - Without parameter

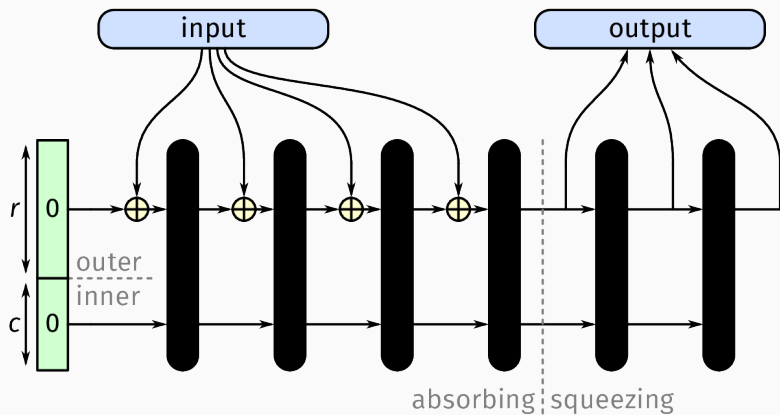
[https://tools.ietf.org/html/  
draft-viguier-kangarootwelve-01](https://tools.ietf.org/html/draft-viguier-kangarootwelve-01)



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# Generic security of the sponge construction

**Theorem 2.** *A padded sponge construction calling a random permutation,  $\mathcal{S}'[\mathcal{F}]$ , is  $(t_D, t_S, N, \epsilon)$ -indistinguishable from a random oracle, for any  $t_D, t_S = O(N^2)$ ,  $N < 2^c$  and for any  $\epsilon$  with  $\epsilon > f_P(N)$ .*

If  $N$  is significantly smaller than  $2^c$ ,  $f_P(N)$  can be approximated closely by:

$$f_P(N) \approx 1 - e^{-\frac{(1-2^{-r})N^2 + (1+2^{-r})N}{2^{c+1}}} < \frac{(1-2^{-r})N^2 + (1+2^{-r})N}{2^{c+1}}. \quad (6)$$

[EuroCrypt 2008]

<http://sponge.noekeon.org/SpongeIndifferentiability.pdf>

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## Theorem, explained

$$\Pr[\text{attack}] \leq \frac{N^2}{2^{c+1}} \text{ (or so)}$$

⇒ if  $N \ll 2^{c/2}$ , then the probability is negligible

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  - Strong mathematical proofs

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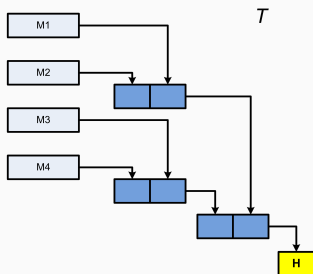
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    - ⇒ lots of third-party **cryptanalysis!**
  - Confidence
    - ⇐ sustained cryptanalysis activity and no break
    - ⇐ proven properties

# Impact of parallelism

$\text{KECCAK-}f[1600] \times 1$	1070 cycles
$\text{KECCAK-}f[1600] \times 2$	1360 cycles
$\text{KECCAK-}f[1600] \times 4$	1410 cycles

CPU: Intel® Core™ i5-6500 (Skylake) with AVX2 256-bit SIMD

# Tree hashing



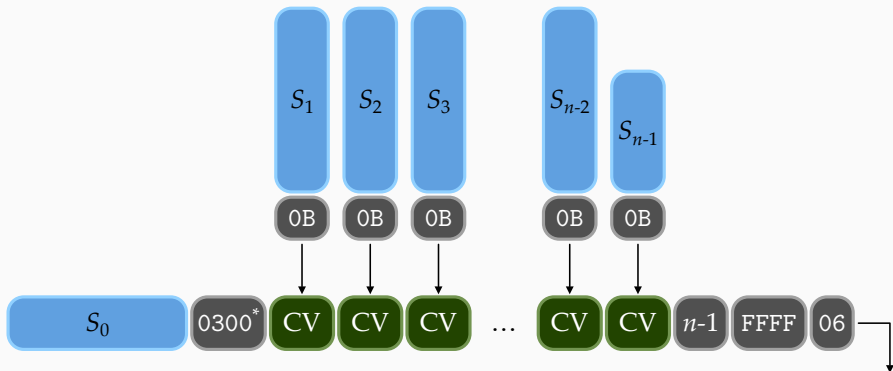
Example: **ParallelHash** [SP 800-185]

function	instruction set	cycles/byte <sup>1</sup>
$\text{KECCAK}[c = 256] \times 1$	x86_64	6.29
$\text{KECCAK}[c = 256] \times 2$	AVX2	4.32
$\text{KECCAK}[c = 256] \times 4$	AVX2	2.31

CPU: Intel® Core™ i5-6500 (Skylake) with AVX2 256-bit SIMD

<sup>1</sup>for long messages.

# KANGAROOTWELVE's mode



Final node growing with kangaroo hopping and SAKURA coding

[ACNS 2014]